

# Architecture for Error Correction: Majority Vote vs. Repeated Parity (Surface Code)

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UCSB and Google

Figure of Merit:  
Threshold  
Qubit number

(Difference is double-exponential)

# Majority Vote vs. Repeated Parity

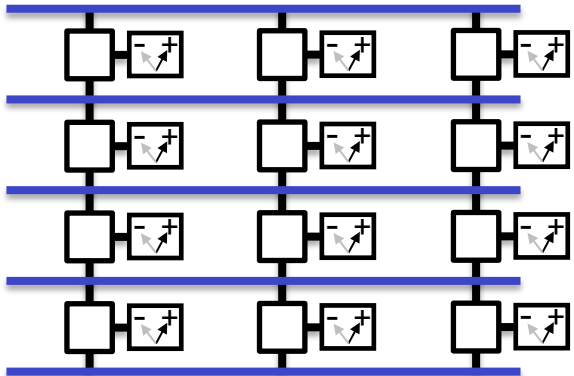
Majority  
Vote



Simple  
Data & measure errors same  
No threshold

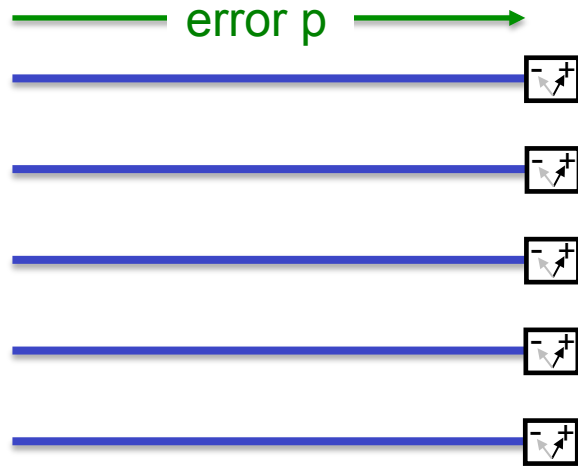
↑ Logical:  
00000 or  
↓ 11111

Repeated  
Parity



Better resource scaling

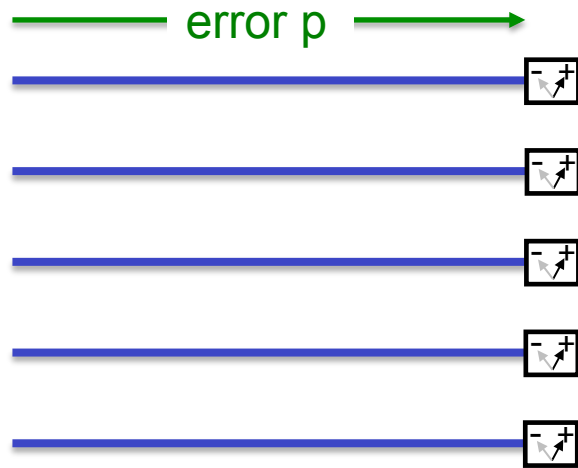
# Logical Error for Majority Vote



$$P_l = \sum_{j=n/2}^n \binom{n}{j} p^j (1-p)^{n-j}$$
$$\cong 0.5 [4p(1-p)]^{n/2}$$
$$\rightarrow 0 \text{ for } n \rightarrow \infty$$

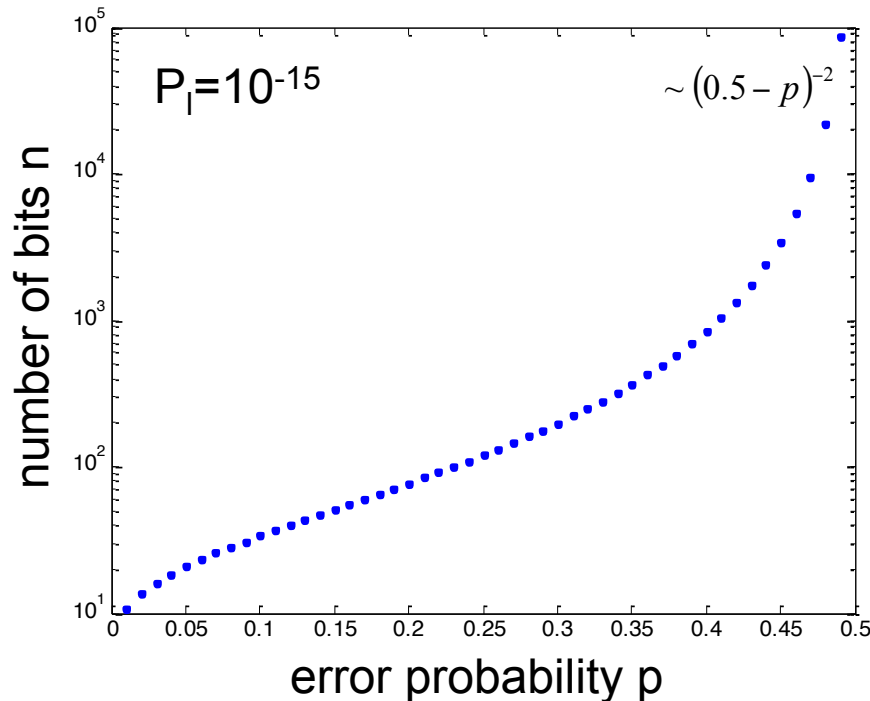
Any information ( $p \neq 0.5$ )  
gives subthreshold

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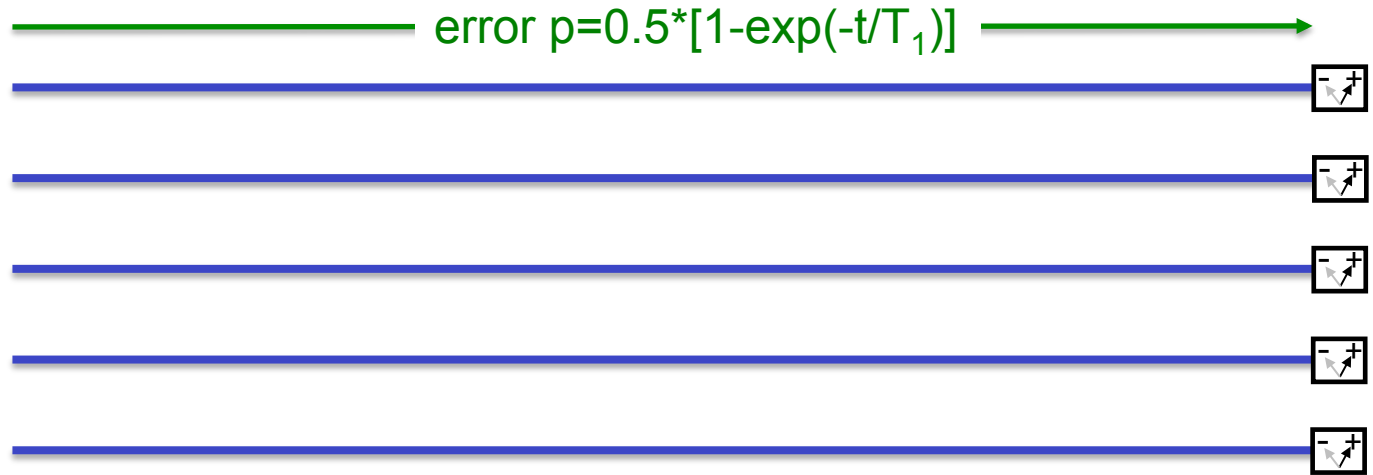
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gives subthreshold



But, number of bits impractical  
as  $p \rightarrow 0.5$

# Error of Majority Vote for Long Time

(symmetrize error)



$$n \approx \exp(2t / T_1)$$

Bad to require exponentially large resources

# Resource Scaling (CS figure of merit)

x is scaling parameter

examples: { Google: number of users  $\sim 10^9$   
Error correction: time  $x=2t/T_1 \sim 10^9$

## resource scaling

1

hashing

$\log x$

$x^{1/2}$

x

$x \log x$

FFT, quicksort

$x^2$

FT

$x^3$

...

majority vote

$\exp x$

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scalable  
↑

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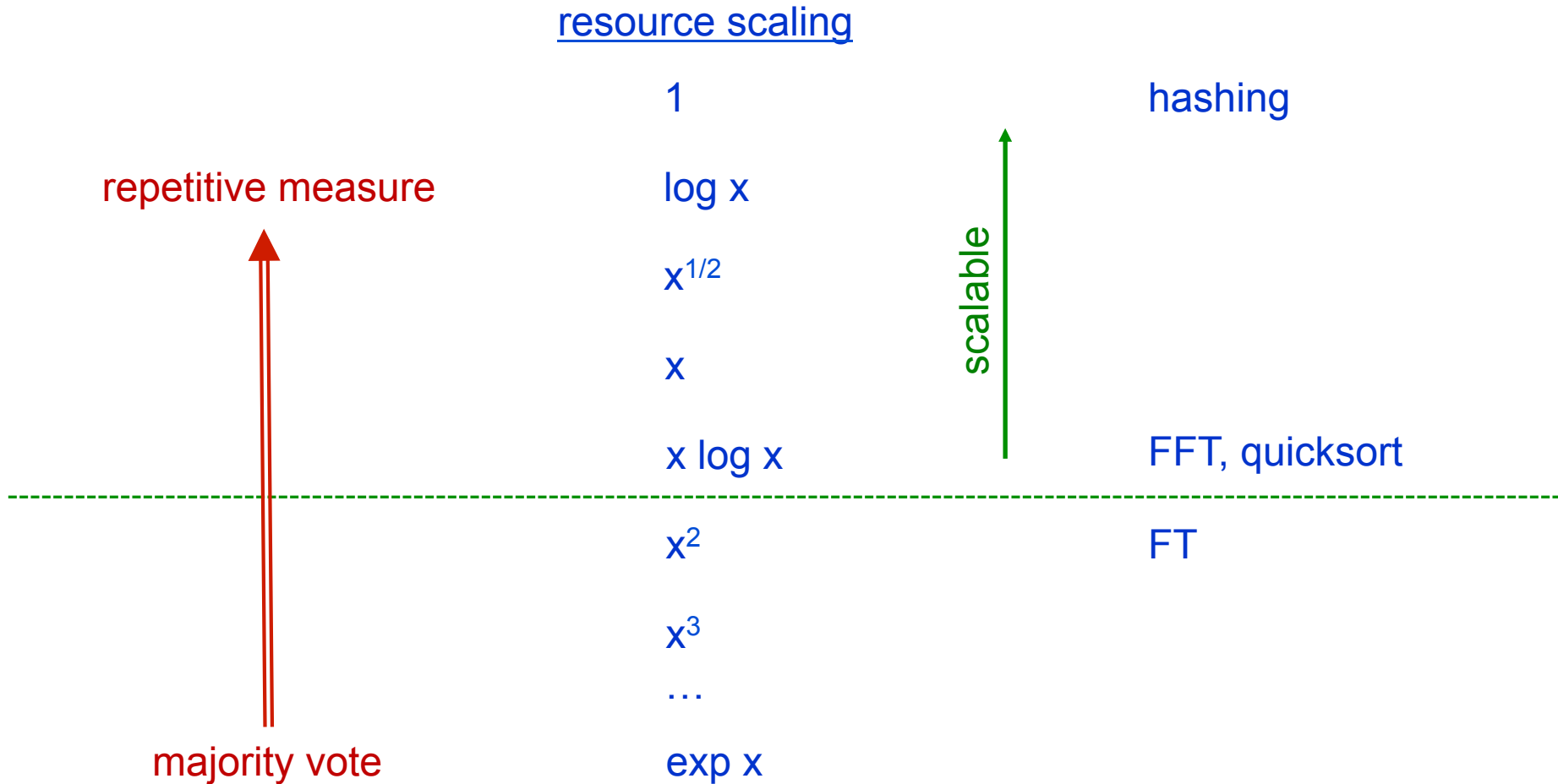
majority vote

$\exp x$

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# Repeated Parity (Surface Code)

0000... & 1111...  
gives same parity

- 1) Repeated measure: log x scaling
- 2) Parity measurement is compatible with quantum and surface code

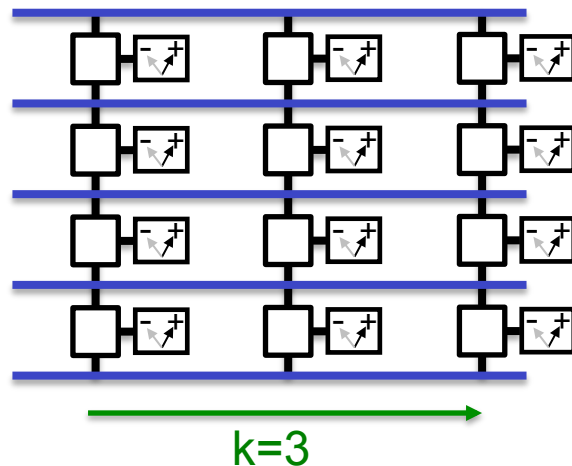
More complex, more error sites

For good scaling, consider only small data errors

$$P_{lk} = 0.5k(4p)^{n/2}$$

For 10x each order  $n/2$ ,  
want  $p = 0.025$

$n/2$  scales as  $\log k$



Repeated  
Parity

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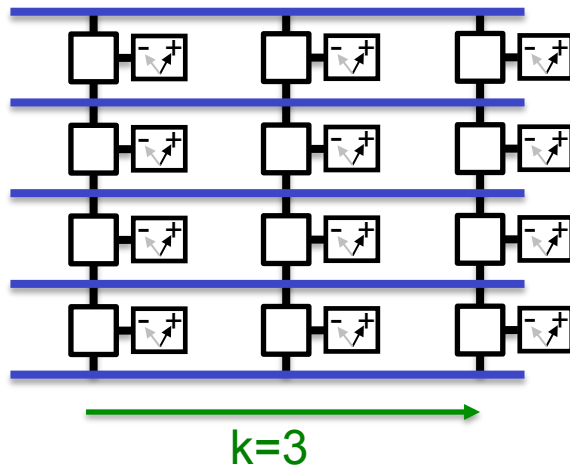
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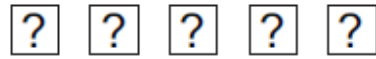
For majority vote, same time

$$P_l = 0.5(4kp)^{n/2}$$

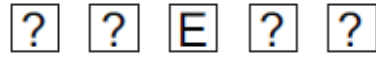
Repeated  
Parity



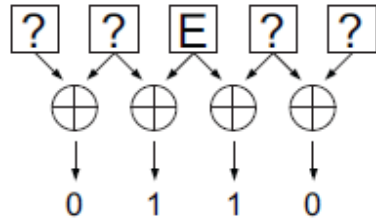
# Bit-Flip Error Correction of Data - Decoding



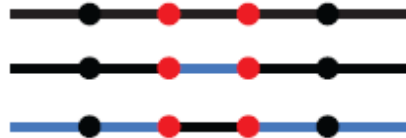
? = 0 or 1



E = bit flip error (with probability  $p$ )



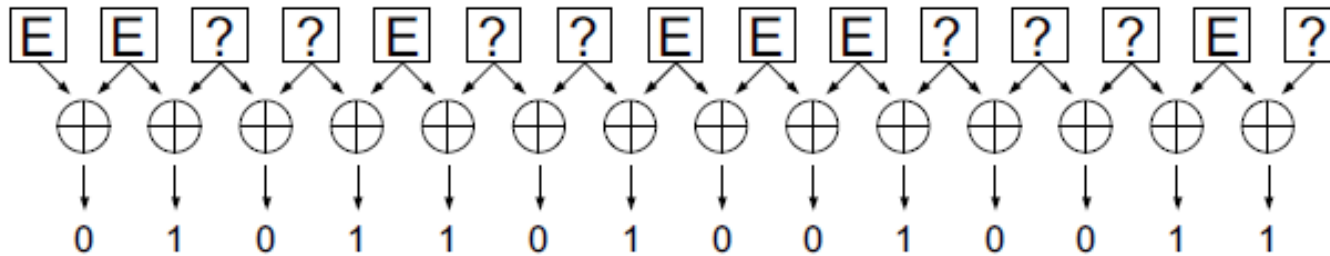
XOR measurement of error



Red dot = error detected

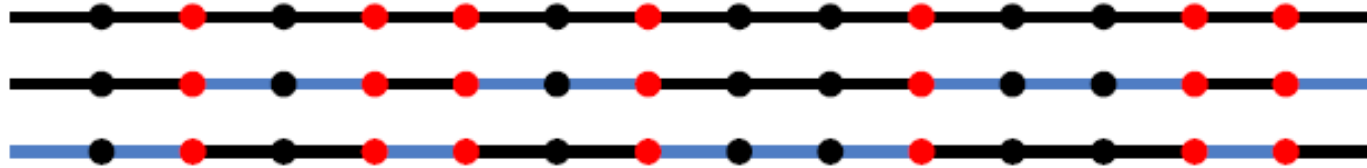
Blue line = error decoded into bit flip

Also possible, but less probable ( $p^4$  vs  $p$ )



Logical error in decoding with probability

$$P_L \cong \frac{15!}{8!7!} p^8$$

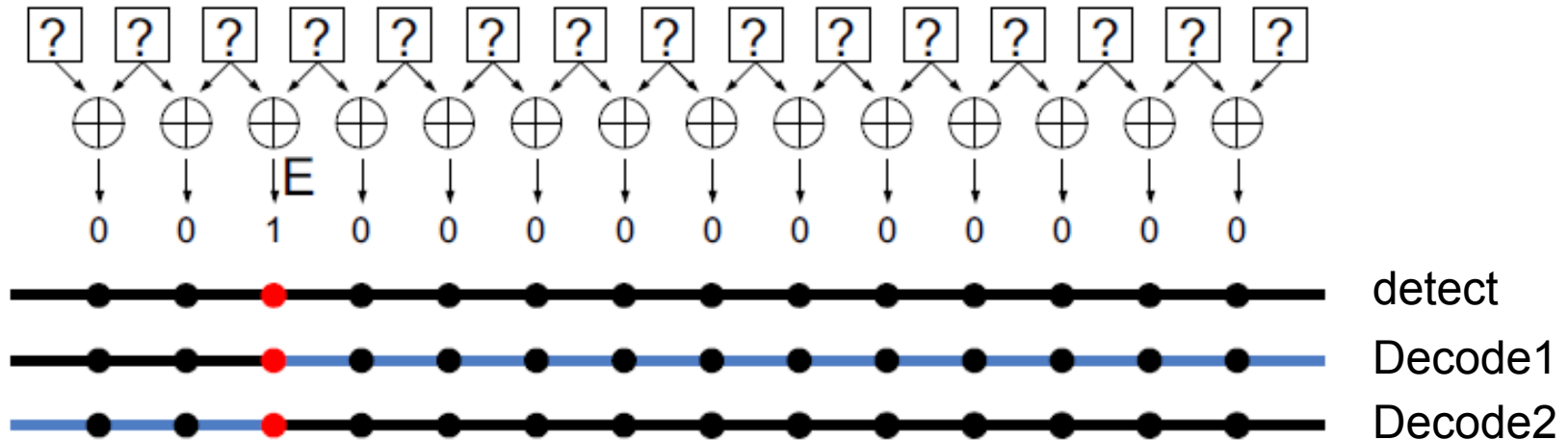


$p^8$

$p^7$

more likely

# Subtlety is for Measurement Errors!

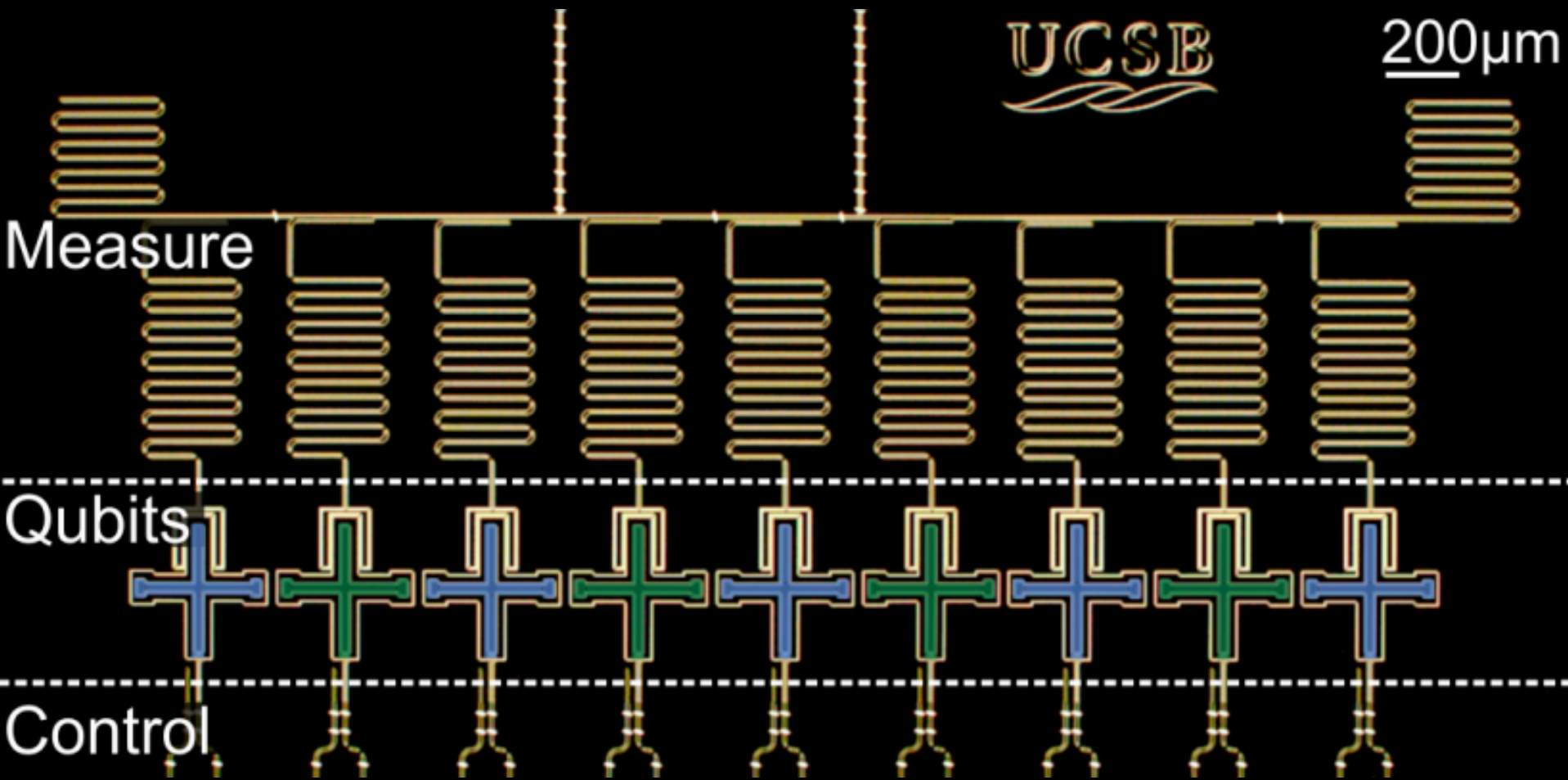


Both decodings are wrong!!

Need to look at measurement vs. time

(The difficulty of quantum hardware)

# Repetition Code: device



# 9 Qubit Experiment: Example data

Error **detection** and **decoding**:

in 000000000

0 0 0 0

0 0 0 0

1 0 1 1 0

0 1 1 0

. D D .

2 0 0 0 0

0 0 0 0

. . . .

3 0 1 1 0

0 1 1 0

. . . .

4 0 0 0 0

0 0 0 0

. . . .

5 1 0 1 0

1 0 1 0

D D . .

6 0 0 0 0

0 0 0 0

. . . .

7 1 0 0 0

1 0 0 0

. . D .

8 0 0 0 0

0 0 0 0

. . . .

0 1 1 0

D D D .

fin 0 0 1 0 0

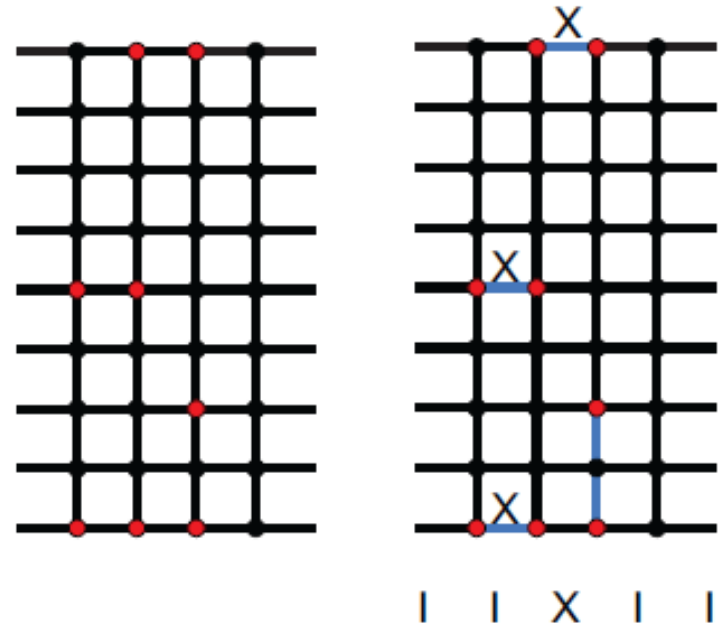
raw data

extended

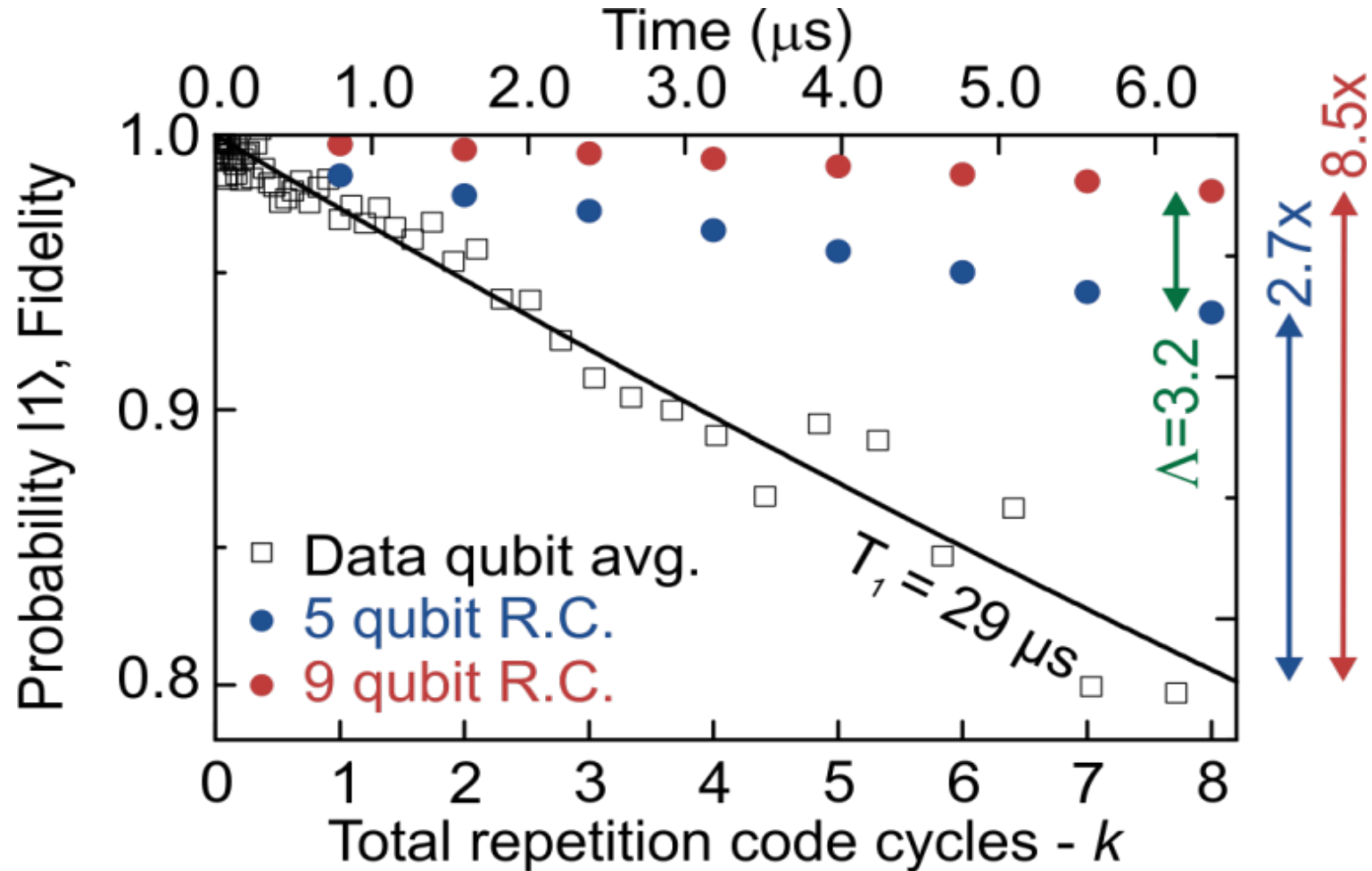
detected  
errors

graph

decoding



# 9 Qubit Data: Bit-Flip Error Correction Works!



Slower decay of logical qubit state:

1<sup>st</sup> order: 2.7x

2<sup>nd</sup> order: 8.5x

$\Lambda = 3.2 > 1$ , so better memory for higher order (fault tolerant behavior)





# Summary and Conclusions

- 1) Impact/meaning of “below threshold” depends on resource scaling
  - majority:  $\exp(10^9)$
  - repetitive:  $\log(10^9)$
- 2) Need to make gates better